Simplification Goals

The goal of the simplification is to minimize the cost of realizing a function with physical circuit elements.

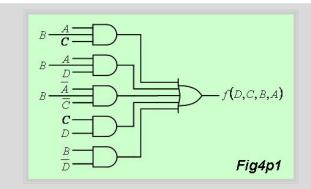
Save silicon area in VLSI circuits or reduce the number of SSI-circuits in implementation.

In general, it is desirable to minimize the number of circuit elements and to make each element as simple as possible.

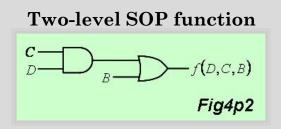
Reduce the number of product terms to reduce gate count.

Minimize the number of literals in each product terms to minimize the number of gate inputs.

> SSI 1-10 MSI 10-100 LSI 100-10 000 VLSI >100 000



Unsimplified two-level switching function has 4 variables, 13 literals and 5 product terms.



Simplified switching function has 3 literals and 2 product terms (CD, B) and 3 variables.

Simplification

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Three "paper and pencil" methods

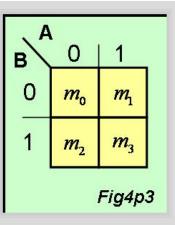
Karnaugh Map method up to six variables

Quine-McCluskey Tabular Minimization method

Switching algebra

Simplification

Karnaugh map for variables B and A



All possible minterms for variables B and A are : m_0 , m_1 , m_2 and m_3 .

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All possible maxterms for variables B and A are : M_0 , M_1 , M_2 and M_3 .

$$m_{0} = \overline{B} \cdot \overline{A} \qquad \qquad M_{0} = B + A$$

$$m_{1} = \overline{B} \cdot A \qquad \qquad M_{1} = B + \overline{A}$$

$$m_{2} = B \cdot \overline{A} \qquad \qquad M_{2} = \overline{B} + A$$

$$m_{3} = B \cdot A \qquad \qquad M_{3} = \overline{B} + \overline{A}$$

3

Simplification

Karnaugh map for variables C, B and A

All possible minterms for variables C, B and A are $: m_0$, m_1, m_2, \dots and m_7 .

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All possible maxterms for variables C, B and A are : M_0 , M_1 , M_2 , and M_7 .

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A

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A

$$m_{0} = \overline{C} \cdot \overline{B} \cdot \overline{A} \qquad \qquad M_{0} = C + B + A$$

$$m_{1} = \overline{C} \cdot \overline{B} \cdot A \qquad \qquad M_{1} = C + B + \overline{A}$$

$$m_{2} = \overline{C} \cdot B \cdot \overline{A} \qquad \qquad M_{2} = C + \overline{B} + A$$

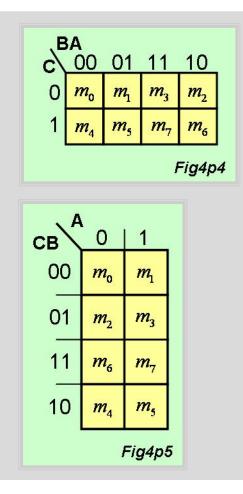
$$m_{3} = \overline{C} \cdot B \cdot A \qquad \qquad M_{3} = C + \overline{B} + \overline{A}$$

$$m_{4} = C \cdot \overline{B} \cdot \overline{A} \qquad \qquad M_{4} = \overline{C} + B + A$$

$$m_{5} = C \cdot \overline{B} \cdot A \qquad \qquad M_{5} = \overline{C} + B + \overline{A}$$

$$m_{6} = C \cdot B \cdot \overline{A} \qquad \qquad M_{6} = \overline{C} + \overline{B} + A$$

$$m_{7} = \overline{C} + \overline{B} + \overline{A}$$



Simplification

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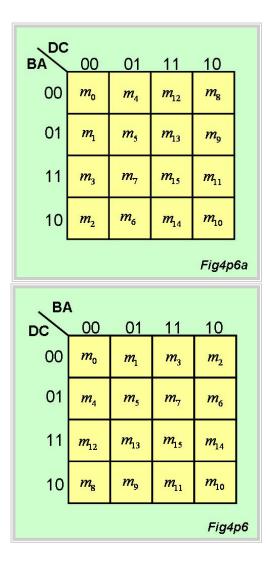
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Karnaugh map for variables D, C, B and A

All possible minterms for variables D, C, B and A .	All possible maxterms for variables D, C, B and A .
-	-
$m_{10} = D \cdot C \cdot B \cdot A$ $m_{11} = D \cdot \overline{C} \cdot B \cdot A$ $m_{12} = D \cdot C \cdot \overline{B} \cdot \overline{A}$ $m_{13} = D \cdot C \cdot \overline{B} \cdot A$ $m_{14} = D \cdot C \cdot B \cdot \overline{A}$ $m_{15} = D \cdot C \cdot B \cdot A$	$\begin{split} M_{10} &= D + C + B + A \\ M_{11} &= \overline{D} + C + \overline{B} + \overline{A} \\ M_{12} &= \overline{D} + \overline{C} + B + A \\ M_{13} &= \overline{D} + \overline{C} + B + \overline{A} \\ M_{14} &= \overline{D} + \overline{C} + \overline{B} + A \\ M_{15} &= \overline{D} + \overline{C} + \overline{B} + \overline{A} \end{split}$



Simplification

Karnaugh map for variables E, D, C, B and A

K-map for variables E , D , C , B and A.

DC BA	00	01	11	10	, DC	00	01	11	10	,
00	m _o	m ₁	<i>m</i> ₃	<i>m</i> ₂	00	<i>m</i> 16	<i>m</i> ₁₇	<i>m</i> 19	m ₁₈	
01	<i>m</i> ₄	<i>m</i> 55	m ₇	<i>m</i> ₆	01	<i>m</i> ₂₀	<i>m</i> ₂₁	<i>m</i> ₂₃	<i>m</i> ₂₂	
11	<i>m</i> ₁₂	<i>m</i> ₁₃	<i>m</i> ₁₅	<i>m</i> ₁₄	11	<i>m</i> ₂₈	<i>m</i> ₂₉	<i>m</i> ₃₁	<i>m</i> ₃₀	
10	m ₈	<i>m</i> 99	<i>m</i> ₁₁	<i>m</i> ₁₀	10	<i>m</i> 24	<i>m</i> ₂₅	<i>m</i> ₂₇	m ₂₆	
		E=(0				E	=1	Fig4p	o7

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Alternative K-map format for five variables

ED	с				、ED	с			
ВА	000	001	011	010	ВА	100	101	111	110
00	m _o	<i>m</i> ₄	<i>m</i> ₁₂	m ₈	00	<i>m</i> ₁₆	<i>m</i> ₂₀	<i>m</i> ₂₈	<i>m</i> ₂₄
01	m	<i>m</i> 5	<i>m</i> ₁₃	<i>m</i> 99	01	<i>m</i> ₁₇	<i>m</i> ₂₁	m ₂₉	<i>m</i> ₂₅
11	m ₃	<i>m</i> 77	<i>m</i> ₁₅	<i>m</i> ₁₁	11	<i>m</i> 19	<i>m</i> ₂₃	<i>m</i> ₃₁	<i>m</i> ₂₇
10	m ₂	<i>m</i> ₆	<i>m</i> ₁₄	<i>m</i> ₁₀	10	m ₁₈	<i>m</i> ₂₂	<i>m</i> ₃₀	<i>m</i> ₂₆
									Fig4p9

Simplification

Karnaugh map for variables E, D, C, B and A

All possible minterms for variables E, D, C, B and A are : m_0 , m_1 , m_2 , m_3 , and m_{31} .

$$\begin{split} m_{0} &= \overline{E} \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot \overline{A} \quad m_{4} = \overline{E} \cdot \overline{D} \cdot C \cdot \overline{B} \cdot \overline{A} \quad m_{8} = \overline{E} \cdot D \cdot \overline{C} \cdot \overline{B} \cdot \overline{A} \quad m_{17} = E \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot A \\ m_{1} &= \overline{E} \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot A \quad m_{5} = \overline{E} \cdot \overline{D} \cdot C \cdot \overline{B} \cdot A \quad m_{9} = \overline{E} \cdot D \cdot \overline{C} \cdot \overline{B} \cdot A \quad m_{18} = E \cdot \overline{D} \cdot \overline{C} \cdot B \cdot \overline{A} \\ m_{2} &= \overline{E} \cdot D \cdot \overline{C} \cdot B \cdot \overline{A} \quad m_{6} = \overline{E} \cdot \overline{D} \cdot C \cdot B \cdot \overline{A} \quad m_{10} = \overline{E} \cdot D \cdot \overline{C} \cdot B \cdot \overline{A} \quad \cdots \\ m_{3} &= \overline{E} \cdot \overline{D} \cdot \overline{C} \cdot B \cdot A \quad m_{7} = \overline{E} \cdot \overline{D} \cdot C \cdot B \cdot A \quad \cdots \\ \end{split}$$

All possible maxterms for variables E, D, C, B and A are : M_0 , M_1 , M_2 , M_3 , ... and M_{31} .

$$\begin{split} M_{0} &= E + D + C + B + A \quad M_{4} = E + D + \overline{C} + B + A \quad M_{8} = E + \overline{D} + C + B + A \quad M_{17} = \overline{E} + D + C + B + \overline{A} \\ M_{1} &= E + D + C + B + \overline{A} \quad M_{5} = E + D + \overline{C} + B + \overline{A} \quad M_{9} = E + \overline{D} + C + B + \overline{A} \quad M_{18} = \overline{E} + D + C + \overline{B} + A \\ M_{2} &= E + D + C + \overline{B} + A \quad M_{6} = E + D + \overline{C} + \overline{B} + A \quad M_{10} = E + \overline{D} + C + \overline{B} + A \quad \cdots \\ M_{3} &= E + D + C + \overline{B} + \overline{A} \quad M_{7} = E + D + \overline{C} + \overline{B} + \overline{A} \quad \cdots \\ M_{31} &= \overline{E} + \overline{D} + \overline{C} + \overline{B} + \overline{A} \quad \cdots \\ \end{split}$$

Simplification

Karnaugh map for variables F, E, D, C, B and A

DC	00	01	11	10	DC B4	00	01	11	10
00	m _o	m	<i>m</i> ₃	m ₂	00	<i>m</i> ₁₆	<i>m</i> ₁₇	<i>m</i> 19	<i>m</i> ₁₈
01	<i>m</i> ₄	<i>m</i> 5	m ₇	m ₆	01	<i>m</i> ₂₀	<i>m</i> ₂₁	<i>m</i> ₂₃	<i>m</i> ₂₂
11	<i>m</i> ₁₂	<i>m</i> ₁₃	<i>m</i> ₁₅	<i>m</i> ₁₄	11	<i>m</i> ₂₈	<i>m</i> ₂₉	<i>m</i> ₃₁	<i>m</i> ₃₀
10	m ₈	m ₉	<i>m</i> ₁₁	<i>m</i> ₁₀	10	<i>m</i> ₂₄	<i>m</i> ₂₅	<i>m</i> ₂₇	<i>m</i> ₂₆
FE=00					FE=01				
DC B/	00	01	11	10	DC BA	00	01	11	10
00	<i>m</i> ₄₈	<i>m</i> 49	<i>m</i> ₅₁	<i>m</i> ₅₀	00	<i>m</i> ₃₂	<i>m</i> ₃₃	<i>m</i> 35	<i>m</i> ₃₄
01	<i>m</i> ₅₂	<i>m</i> ₅₃	<i>m</i> ₅₅	<i>m</i> 54	01	<i>m</i> ₃₆	m ₃₇	m39	m ₃₈
11	<i>m</i> ₆₀	<i>m</i> ₆₁	m ₆₃	<i>m</i> ₆₂	11	<i>m</i> ₄₄	<i>m</i> ₄₅	<i>m</i> 47	<i>m</i> ₄₆
10	<i>m</i> 556	m ₅₇	m ₅₉	m ₅₈	10	<i>m</i> ₄₀	<i>m</i> ₄₁	m ₄₃	<i>m</i> ₄₂
		FE	=11				F	E=10	Fig4p10

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Simplification

Karnaugh map for variables F, E, D, C, B and A

All possible minterms for variables F, E, D, C, B and A are : m_0 , m_1 , m_2 , m_3 ... and m_{63} .

$$\begin{split} m_{0} &= \overline{F} \cdot \overline{E} \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot \overline{A} \quad m_{16} = \overline{F} \cdot E \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot \overline{A} \quad m_{32} = F \cdot \overline{E} \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot \overline{A} \quad m_{48} = F \cdot E \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot \overline{A} \\ m_{1} &= \overline{F} \cdot \overline{E} \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot A \quad m_{17} = \overline{F} \cdot E \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot A \quad m_{33} = F \cdot \overline{E} \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot A \quad m_{49} = F \cdot E \cdot \overline{D} \cdot \overline{C} \cdot \overline{B} \cdot A \\ \vdots &\vdots &\vdots &\vdots \\ m_{15} &= \overline{F} \cdot \overline{E} \cdot D \cdot C \cdot B \cdot A \quad m_{31} = \overline{F} \cdot E \cdot D \cdot C \cdot B \cdot A \quad m_{47} = F \cdot \overline{E} \cdot D \cdot C \cdot B \cdot A \quad m_{63} = F \cdot E \cdot D \cdot C \cdot B \cdot A \end{split}$$

All possible maxterms for variables F, E, D, C, B and A are : M_0 , M_1 , M_2 , . . and M_{63} .

$$\begin{split} M_{\theta} &= F + E + D + C + B + A \\ M_{16} &= F + \overline{E} + D + C + B + A \\ M_{17} &= F + \overline{E} + D + C + B + \overline{A} \\ \vdots &= F + E + D + C + B + \overline{A} \\ \vdots &= F + E + D + C + B + \overline{A} \\ \vdots &= F + E + \overline{D} + \overline{C} + \overline{B} + \overline{A} \\ M_{17} &= F + \overline{E} + D + C + B + \overline{A} \\ \vdots &= F + E + \overline{D} + \overline{C} + \overline{B} + \overline{A} \\ M_{17} &= F + \overline{E} + \overline{D} + \overline{C} + \overline{B} + \overline{A} \\ M_{21} &= F + \overline{E} + \overline{D} + \overline{C} + \overline{B} + \overline{A} \\ M_{47} &= \overline{F} + E + \overline{D} + \overline{C} + \overline{B} + \overline{A} \\ M_{63} &= \overline{F} + \overline{E} + \overline{D} + \overline{C} + \overline{B} + \overline{A} \\ \end{split}$$

Simplification

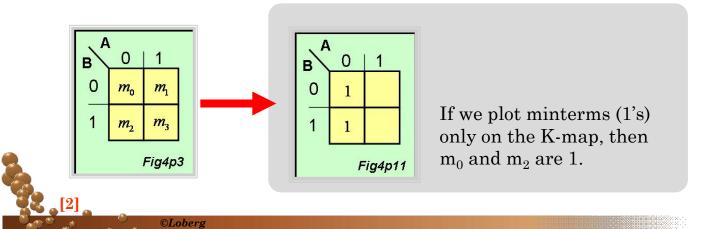
Plotting Functions in Canonical Form on the Karnaugh map

If switching function is expressed in canonical form, it may be readily plotted on a K-map. Each cell on the K-map corresponds to minterm/maxterm of the canonical form.

Suppose we have following Boolean function f_1 which is given in canonical SOP form : $f_1(B, A) = \overline{B}$.

We have assumed that right side of variables (B, A) represent a "least significant bit" and left side (B) represent a "most significant bit". $f_{I}(B,A) = \overline{B} \cdot \overline{A} + B \cdot \overline{A} \quad f_{I}(B,A) = \sum m(0,2)$

⇒	$m_0 = \overline{B} \cdot \overline{A}$
	$m_1 = \overline{B} \cdot A$
	$m_2 = B \cdot \overline{A}$
	$m_3 = B \cdot A$



Simplification

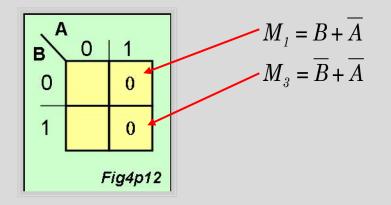
Plotting Functions in Canonical Form on the Karnaugh map

$$f_1(B,A) = \sum m(0,2) = \prod M(1,3) = M_1 \cdot M_3 = (B + \overline{A}) \cdot (\overline{B} + \overline{A})$$

If we plot maxterms (0's) only on the K-map, then M_1 and M_3 are 0.

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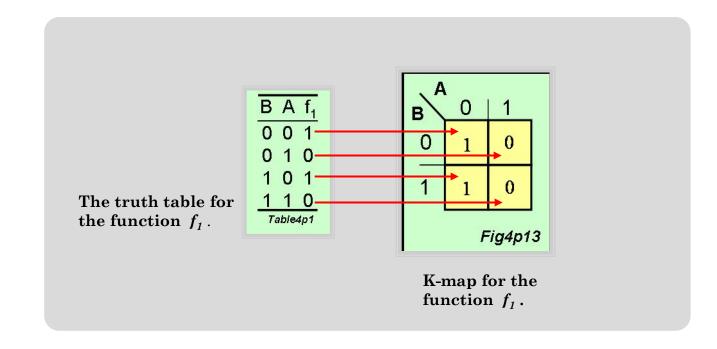


Simplification

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Plotting Functions in Canonical Form on the Karnaugh map

If we plot all minterms (1) and maxterms (0) on the K-map, then the K-map corresponds directly to the truth table of the function f_1 .



Simplification

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$$f_{2}(C, B, A) = \sum m(0,3,5) = m_{0} + m_{3} + m_{5}$$

$$= \overline{C} \cdot \overline{B} \cdot \overline{A} + \overline{C} \cdot B \cdot A + C \cdot \overline{B} \cdot A$$

$$= \prod M(1,2,4,6,7) = M_{1} \cdot M_{2} \cdot M_{4} \cdot M_{6} \cdot M_{7}$$

$$= (C + B + \overline{A}) \cdot (C + \overline{B} + A) \cdot (\overline{C} + B + A) \cdot (\overline{C} + \overline{B} + A) \cdot (\overline{C} + \overline{B} + \overline{A})$$

$$\frac{\overline{C} \ \overline{B} \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ \overline{B} \ A \ f_{2}}{0 \ 0 \ 1 \ 0}$$

$$\frac{\overline{C} \ \overline{B} \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ \overline{B} \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ \overline{B} \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ \overline{B} \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ \overline{B} \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ \overline{B} \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

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$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1}$$

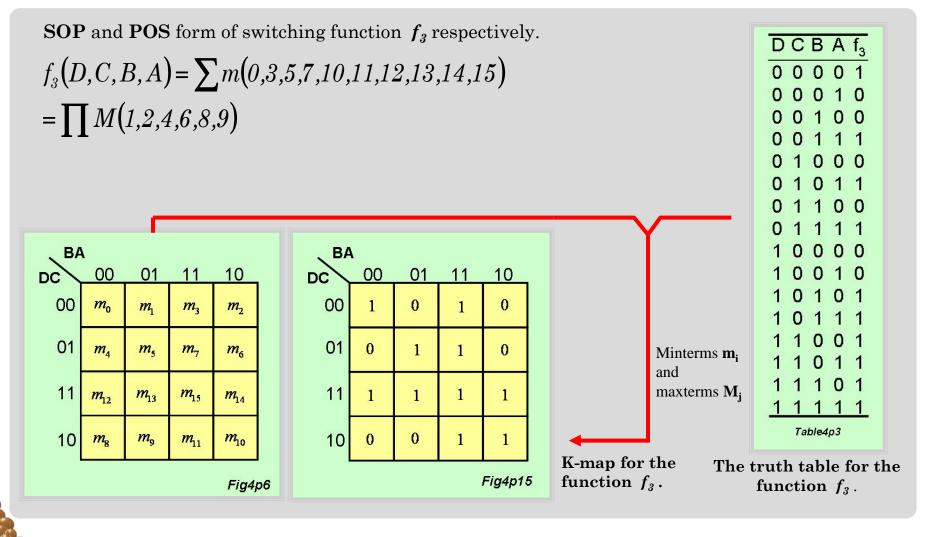
$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1}$$

$$\frac{\overline{C} \ B \ A \ f_{2}}{0 \ 0 \ 1}$$

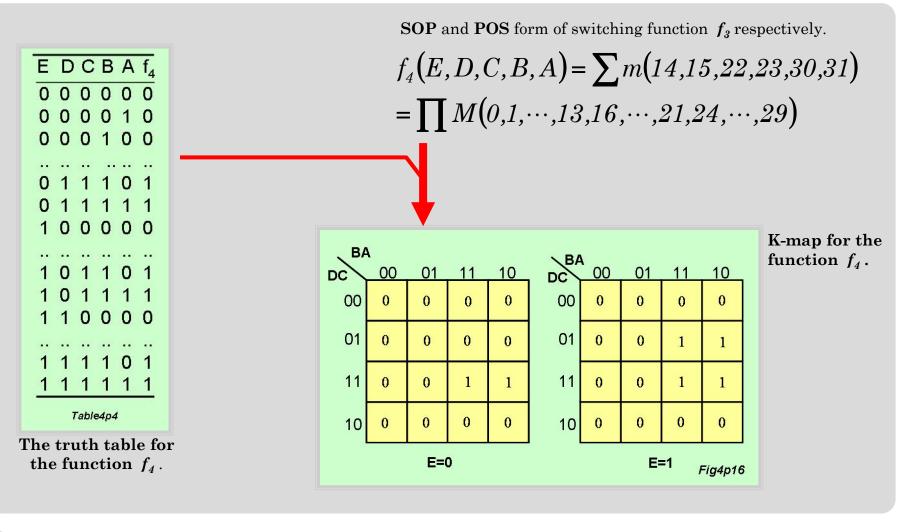
$$\frac{\overline{C} \ B \ A \ f_$$

D, C, B and A

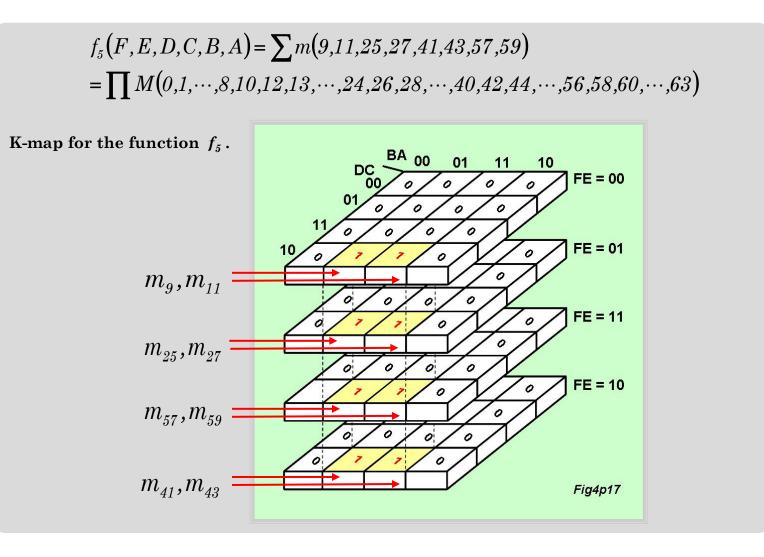
Simplification



Simplification



Simplification



Simplification

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Plotting Functions in Standard SOP/POS Form on the K-map

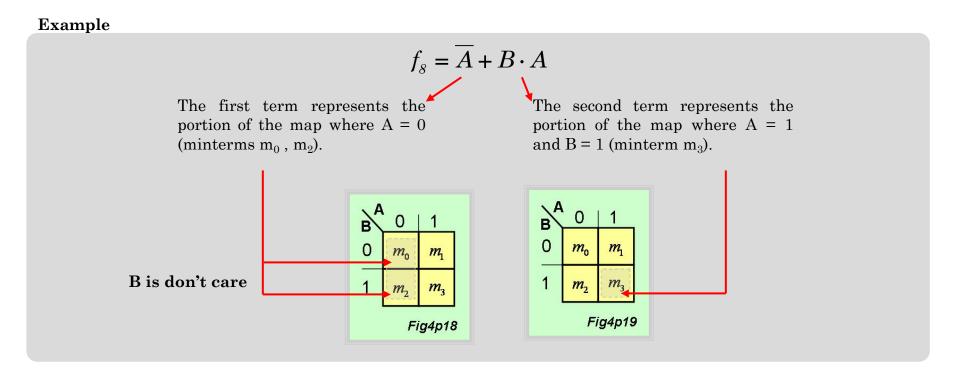
In canonical form the each product or sum term must contain all the variables, either complemented or uncomplemented.

In standard form each	Standard SOP form	$f_6(C, B, A) = A + B \cdot \overline{A} + \overline{C} \cdot B$
product/sum term may have any number of literals.	Standard POS form	$f_7(C, B, A) = (\overline{A} + B + \overline{C}) \cdot (\overline{B} + C) \cdot A$

Simplification

Plotting Functions in Standard SOP/POS Form on the K-map

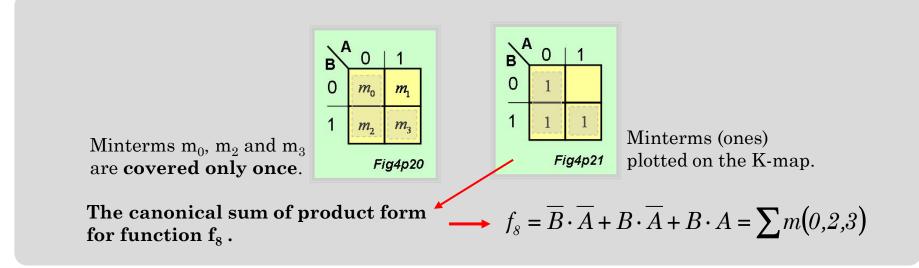
Consider the following function $f_8(B,A)$, which is expressed as a **standard sum of products**.



Variables B and A

Simplification

Plotting Functions in Standard SOP/POS Form on the K-map

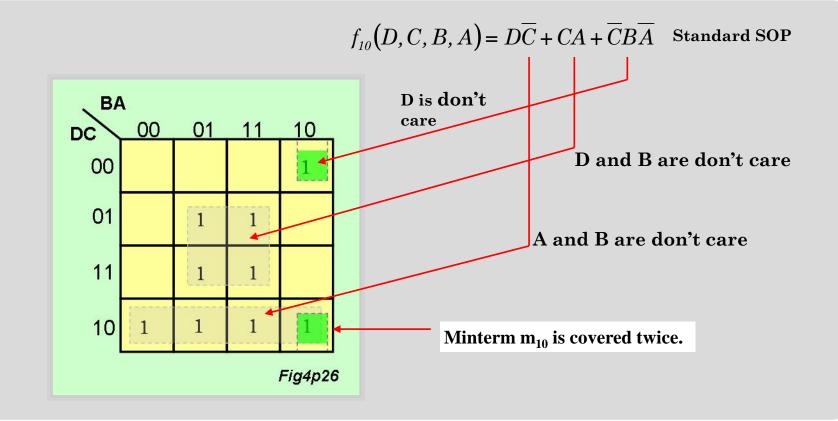


We can use K-map to make a conversion from standard SOP to canonical SOP form.

Simplification Of Switching Functions Variables D, C, B and A

Simplification

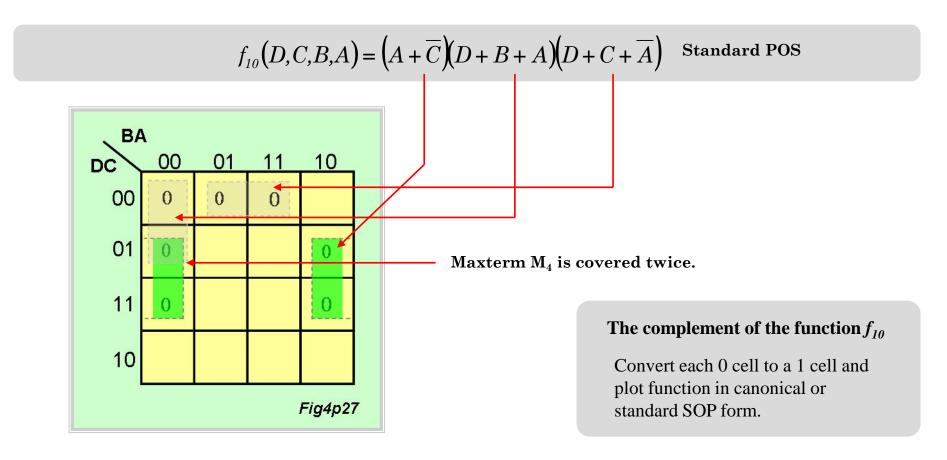
Plotting Functions in Standard SOP/POS Form on the K-map



Simplification Of Switching Functions Variables D, C, B and A

Simplification

Plotting Functions in Standard SOP/POS Form on the K-map



Simplification

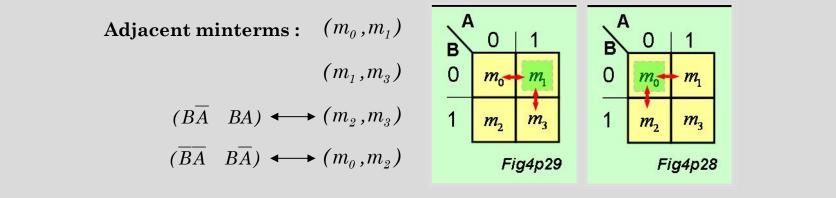
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Guidelines for Simplifying Functions Using K-maps

Five important points to keep in mind.

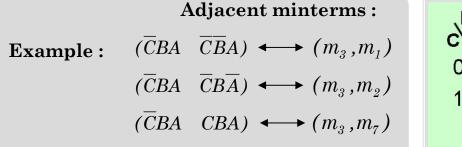
1. Each minterm on a K-map of *n* variables has *n* logically adjacent minterms.

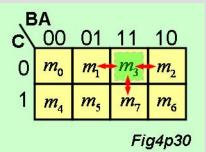


Simplification

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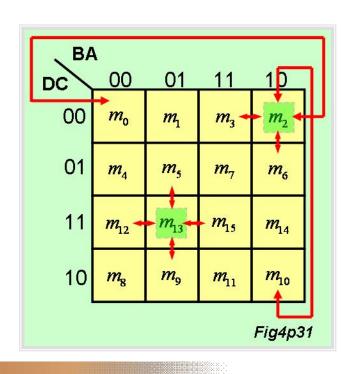
Example :

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 (m_2, m_0) (m_2, m_3) (m_2, m_6) (m_2, m_{10}) **Or** (m_{13}, m_5) (m_{13}, m_9) (m_{13}, m_{12}) (m_{13}, m_{15})



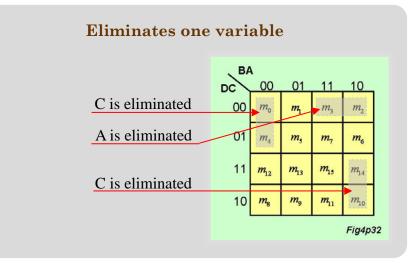
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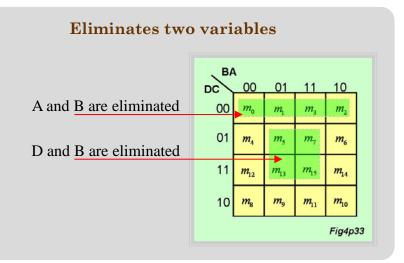
Simplification

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2. Group squares in powers of 2 (2, 4, 8, 16, ..). Grouping 2^n squares, eliminates n variables.





Simplification

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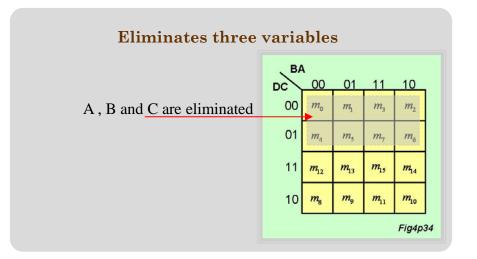
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Simplification

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3. Group as many squares together as possible. The larger the group is, the fewer the number of literals in the resulting product term.

4. Make as few groups as possible to cover all the minterms of function. The fewer the groups, the fewer the number of product terms in the minimized function. Stop, when all minterms are used at least once. The minterm is *covered* when it is used at least once in groups.

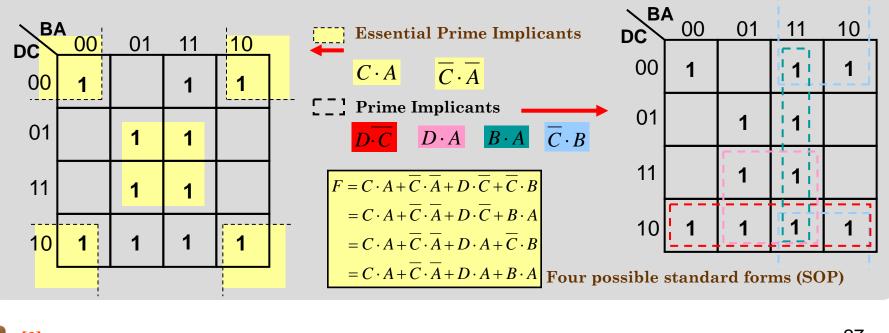
5. Start grouping with those minterms which has fewest number of logically adjacent minterms.

Simplification

General Terminology

We must ensure that all the minterms of the function are covered when we combine squares. There are no minterms already covered by other terms.

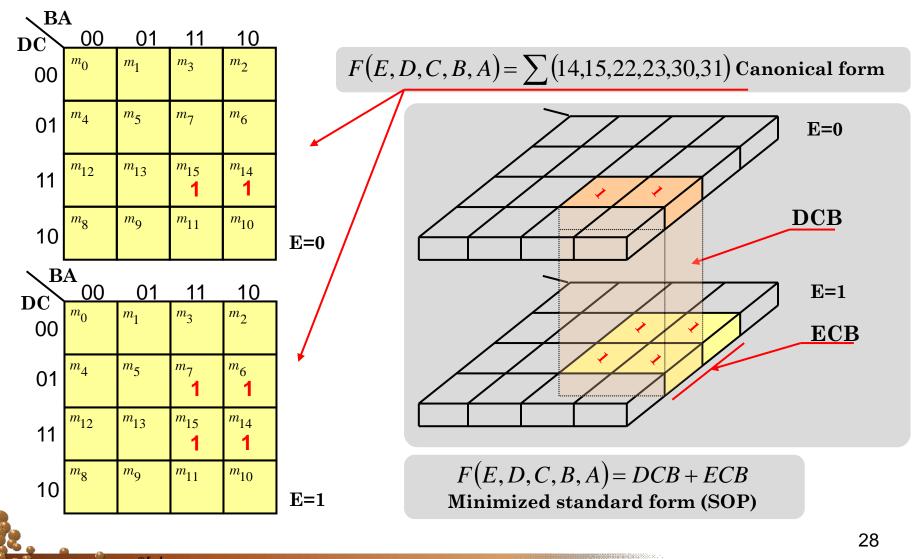
> A *prime implicant* is a product term obtained by combining the maximum possible number of adjacent squares in the map. If *minterm* in a square is *covered by only one prime implicant*, that prime implicant is said to be *essential*.



Simplification

Examples with K-maps

Example with Five-Variable Map

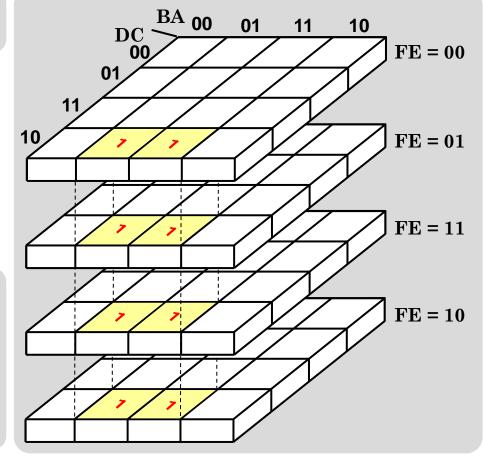


Simplification

Examples with K-maps

Example with Six-Variable Map

 $f_5(F, E, D, C, B, A) = \sum m(9, 11, 25, 27, 41, 43, 57, 59)$



Simplified Boolean function f_5 :

F = DCA

Minimized standard form (SOP)

Simplification

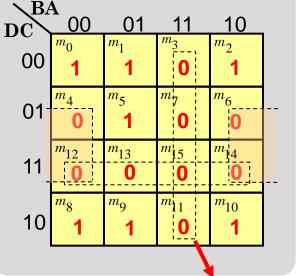
Examples with K-maps

$$F(D,C,B,A) = \sum m(0,1,2,5,8,9,10)$$
$$\prod M(3,4,6,7,11,12,13,14,15)$$

1's marked in the squares (minterms): m_0, m_1, m_2, m_5 ,

 m_8, m_9, m_{10}

0's marked in the squares (maxterms) : $M_{\scriptscriptstyle 3}, M_{\scriptscriptstyle 4}, M_{\scriptscriptstyle 6}, M_{\scriptscriptstyle 7}, M_{\scriptscriptstyle 11},$ $M_{\scriptscriptstyle 12}, M_{\scriptscriptstyle 13}, M_{\scriptscriptstyle 14}, M_{\scriptscriptstyle 15}$



If the squares marked with 0's are combined, (and take reduced minterms) we obtain the simplified complemented function: \overline{F}

If we take reduced maxterms,we obtain simplified function F.

$$F(D,C,B,A) = B \cdot A + D \cdot C + C \cdot A$$

Applying DeMorgan's theorem we obtain simplified F in POS form.

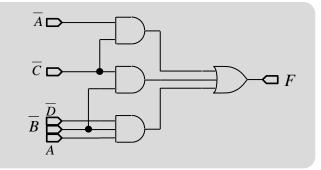
$$F = \overline{\overline{F}}(D, C, B, A) = \overline{B \cdot A} + D \cdot C + C \cdot \overline{A}$$
$$= \overline{(B \cdot A)} \cdot \overline{(D \cdot C)} \cdot \overline{(C \cdot \overline{A})}$$
$$= \overline{(B + \overline{A})} \cdot \overline{(D + \overline{C})} \cdot \overline{(\overline{C} + A)}$$

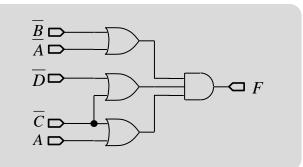
Simplification

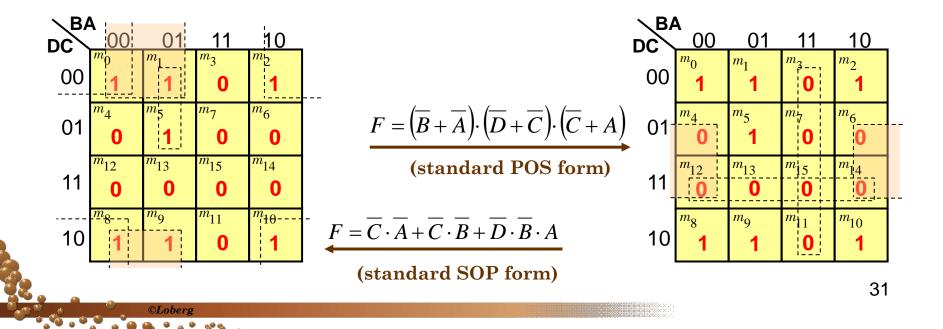
Examples with K-maps

Gate implementation of the function

$$F(D, C, B, A) = \sum (0, 1, 2, 5, 8, 9, 10)$$







Simplification

Examples with K-maps

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In practice, there are some applications where the function is not specified for certain combinations of the input variables.

For example, four-bit binary code for the decimal digits $(0-9, P_0 - P_9)$.

Six unspecified outputs for some input $(P_{10} - P_{15})$ combinations.

These don't care conditions can be used on a map to provide further simplification of the Boolean function.

Mark the don't care minterm with X in the K-map.

$$F = \overline{C} \cdot \overline{B} + \overline{D} \cdot \overline{B} \cdot A + D \cdot \overline{C} \cdot \overline{A} \quad \text{(Simplified without X)}$$

$$F = \overline{C} \cdot \overline{B} + \overline{D} \cdot \overline{B} \cdot A + D \cdot \overline{C} \cdot \overline{A} \quad \text{(Simplified without X)}$$

$$(\text{Simplified with X}) \quad F = \overline{C} \cdot \overline{B} + \overline{D} \cdot \overline{B} \cdot A + \overline{C} \cdot \overline{A}$$

$$(\text{Simplified with X}) \quad F = \overline{C} \cdot \overline{B} + \overline{D} \cdot \overline{B} \cdot A + \overline{C} \cdot \overline{A}$$

The End

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